

• how to make recursive functions more efficient with tail recursion

This week:

- code reuse with higher-order functions (HOFs)
- some useful HOFs: map, filter, and fold

Recursion is good...

- Recursive code mirrors recursive data
 - Base constructor -> Base case
 - Inductive constructor -> Inductive case (with recursive call)
- But it can get kinda repetitive!

Example: evens

Let's write a function evens :

```
-- evens [] ==> []

-- evens [1,2,3,4] ==> [2,4]

evens :: [Int] -> [Int]

evens [] = ...

evens (x:xs) = ...
```

Example: four-letter words

Let's write a function fourChars :

```
-- fourChars [] ==> []
-- fourChars ["i", "must", "do", "work"] ==> ["must", "work"]
fourChars :: [String] -> [String]
fourChars [] = ...
fourChars (x:xs) = ...
```

Yikes, Most Code is the Same

Lets rename the functions to foo :

```
foo [] = []
foo (x:xs)
    | x mod 2 == 0 = x : foo xs
    | otherwise = foo xs
foo [] = []
foo (x:xs)
    | length x == 4 = x : foo xs
    | otherwise = foo xs
```

Only difference is condition

• x mod 2 == 0 vs length x == 4

Moral of the day

D.R.Y. Don't Repeat Yourself!

Can we

-

- *reuse* the general pattern and
- *substitute in* the custom condition?

>





General Pattern

- expressed as a higher-order function
- takes customizable operations as arguments

Specific Operation

• passed in as an argument to the HOF



The filter Pattern

General Pattern

- HOF filter
- Recursively traverse list and pick out elements that satisfy a predicate

Specific Operations

• Predicates isEven and isFour

 $\lambda x. x+1$



filter instances

Avoid duplicating code!

Let's talk about types

```
-- evens [1,2,3,4] ==> [2,4]

evens :: [Int] -> [Int]

evens xs = filter isEven xs

where

isEven :: Int -> Bool

isEven x = x `mod` 2 == 0
```

filter :: ???

```
-- fourChars ["i", "must", "do", "work"] ==> ["must", "work"]
fourChars :: [String] -> [String]
fourChars xs = filter isFour xs
  where
    isFour :: String -> Bool
    isFour x = length x == 4
filter :: ???
```

```
So what's the type of filter?
```

```
filter :: (Int -> Bool) -> [Int] -> [Int] -- ???
```

```
filter :: (String -> Bool) -> [String] -> [String] -- ???
```

- It does not care what the list elements are
 - as long as the predicate can handle them
- It's type is polymorphic (generic) in the type of list elements

```
-- For any type `a`
```

- -- if you give me a predicate on `a`s
- -- and a list of `a`s,

```
-- I'll give you back a list of `a`s
```

filter :: (a -> Bool) -> [a] -> [a]

Example: all caps

Lets write a function shout :

-- shout [] ==> [] -- shout ['h','e','l','l','o'] ==> ['H','E','L','L','O']

```
shout :: [Char] -> [Char]
shout [] = ...
shout (x:xs) = ...
```

Example: squares

Lets write a function squares :

```
-- squares [] ==> []
-- squares [1,2,3,4] ==> [1,4,9,16]
squares :: [Int] -> [Int]
squares [] = ...
squares (x:xs) = ...
```

Yikes, Most Code is the Same

Lets rename the functions to foo :

```
-- shout

foo [] = []

foo (x:xs) = toUpper x : foo xs

-- squares

foo [] = []

foo (x:xs) = (x * x) : foo xs
```

Lets refactor into the common pattern

pattern = ...

The "map" pattern

shout	[]	=	[]				squares	[]	=	[]		
shout	(x:xs)	=	toUpper x	:	shout	xs	squares	(x:xs)	=	(x*x) :	squares	xs

The map Pattern

General Pattern

- HOF map
- Apply a transformation f to each element of a list

Specific Operations

• Transformations to Upper and $x \rightarrow x * x$

map f [] = []
map f (x:xs) = f x : map f xs

Lets refactor shout and squares

shout = map ...

squares = map ...

shout = map (
$$x \rightarrow toUpper x$$
)

squares = map ($x \rightarrow x*x$)

map instances

QUIZ

What is the type of map?

map f [] = []
map f (x:xs) = f x : map f xs
(A) (Char -> Char) -> [Char] -> [Char] X
(B) (Int -> Int) -> [Int] -> [Int] X
(C) (a -> a) -> [a] -> [a] 4
(D) (a -> b) -> [a] -> [b] 4
(E) (a -> b) -> [c] -> [d]

```
-- For any types `a` and `b`
-- if you give me a transformation from `a` to `b`
-- and a list of `a`s,
-- I'll give you back a list of `b`s
map :: (a -> b) -> [a] -> [b]
```

Type says it all!

- The only meaningful thing a function of this type can do is apply its first argument to elements of the list
- Hoogle it!

Things to try at home:

- can you write a function map' :: (a -> b) -> [a] -> [b] whose behavior is different from map?
- can you write a function map' :: (a -> b) -> [a] -> [b] such that map' f xs returns a list whose elements are not in map f xs?

QUIZ

What is the value of quiz ?

map ::
$$(a \rightarrow b) \rightarrow [a] \rightarrow [b]$$

 $(lut, lut) \rightarrow lut$
quiz = map $((x, y) \rightarrow x + y) [1, 2, 3]$
(A) [2, 4, 6]
(B) [3, 5]
(C) Syntax Error

(D) Type Error

(E) None of the above

Don't Repeat Yourself

Benefits of **factoring** code with HOFs:

- Reuse iteration pattern
 - think in terms of standard patterns
 - less to write
 - easier to communicate
- Avoid bugs due to repetition



Recall: length of a list

```
-- len [] ==> 0

-- len ["carne", "asada"] ==> 2

len :: [a] -> Int

len [] = 0

len (x:xs) = 1 + len xs
```

Recall: summing a list

-- sum [] ==> 0 -- sum [1,2,3] ==> 6 sum :: [Int] -> Int sum [] = 0 sum (x:xs) = x + sum xs

Example: string concatenation

Let's write a function cat :

```
-- cat [] ==> ""
-- cat ["carne", "asada", "torta"] ==> "carneasadatorta"
cat :: [String] -> String
cat [] = ...
cat (x:xs) = ...
```

Can you spot the pattern?

```
-- len

foo [] = 0

foo (x:xs) = 1 + foo xs

-- sum

foo [] = 0

foo (x:xs) = x + foo xs

-- cat

foo [] = ""

foo (x:xs) = x ++ foo xs
```

```
Base
```

pattern = ...

The "fold-right" pattern

len [] = 0	sum [] = 0	cat [] = ""
len (x:xs) = $1 + len xs$	sum(x:xs) = x + sum xs	cat (x:xs) = x ++ sum xs

The foldr Pattern

General Pattern

- Recurse on tail
- Combine result with the head using some binary operation

```
foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)
```

Let's refactor sum, len and cat:

sum = foldr
cat = foldr
len = foldr

Factor the recursion out!

foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)

len = foldr $(x n \rightarrow 1 + n)$ 0

sum = foldr (x n - x + n) 0

cat = foldr (\x s -> x ++ n) ""

foldr instances

You can write it more clearly as

sum = foldr (+) 0

cat = foldr (++) ""

QUIZ

What does this evaluate to?

```
foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)
quiz = foldr (:) [] [1,2,3]
(A) Type error
(B) [1,2,3]
(C) [3,2,1]
(D) [[3],[2],[1]]
(E) [[1],[2],[3]]
```

```
foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)

foldr (:) [] [1,2,3]
==> (:) 1 (foldr (:) [] [2, 3])
==> (:) 1 ((:) 2 (foldr (:) [] [3]))
==> (:) 1 ((:) 2 ((:) 3 (foldr (:) [] [])))
==> (:) 1 ((:) 2 ((:) 3 []))
== 1 : (2 : (3 : []))
== [1,2,3]
```

 $((b \circ p \times_1) \circ p \times_2) \circ p \times_3) \circ p \times_4)$ $(((X_1 \circ p \times_2) \circ p \times_3) \circ p \times_4) \circ p \times_6)$ $((X_1 \circ p \times_2) \circ p \times_3) \circ p \times_4) \circ p \times_6)$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6)$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_2 \circ p \times_3) \circ p \times_4) \circ p \times_6$ $(X_1 \circ p \times_4) \circ p \times_6$ $(Y_1 \circ p \times_6) \circ p \times_6$ $(Y_1 \circ p \times_6) \circ p \times_$

==> f x1 (f x2 (f x3 (f x4 (foldr f b [])))

==> f x1 (f x2 (f x3 (f x4 b)))

Accumulate the values from the right

For example:

```
foldr (+) 0 [1, 2, 3, 4]
==> 1 + (foldr (+) 1 [2, 3, 4])
==> 1 + (2 + (foldr (+) 0 [3, 4]))
==> 1 + (2 + (3 + (foldr (+) 0 [4])))
==> 1 + (2 + (3 + (4 + (foldr (+) 0 []))))
==> 1 + (2 + (3 + (4 + 0)))
```

QUIZ

What is the most general type of foldr?

foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)
(A) (a -> a -> a) -> a -> [a] -> a
(B) (a -> a -> b) -> a -> [a] -> b
(C) (a -> b -> a) -> b -> [a] -> b
(D) (a -> b -> b) -> b -> [a] -> b
(E) (b -> a -> b) -> b -> [a] -> b

Is foldr tail recursive?

What about tail-recursive versions?

Let's write tail-recursive sum!

sumTR :: [Int] -> Int
sumTR = ...

Lets run sumTR to see how it works

```
sumTR [1,2,3]
==> helper 0 [1,2,3]
==> helper 1 [2,3] -- 0 + 1 ==> 1
==> helper 3 [3] -- 1 + 2 ==> 3
==> helper 6 [] -- 3 + 3 ==> 6
==> 6
```

Note: helper directly returns the result of recursive call!