

Recursion \equiv FIXPOINT COMBINATOR

Haskell Crash Course Part I

From the Lambda Calculus to Haskell

What is Haskell?

A typed, lazy, purely functional programming language

Haskell = λ -calculus ++

- better syntax ✓

- types

ONE = $\lambda f x \rightarrow f x$

TWO = $\lambda f x \rightarrow f (f x)$

(ONE TWO)

- built-in features

- booleans, numbers, characters
- records (tuples)
- lists
- recursion
- ...

(1 2)

↑ TYPE ERROR

Programming in Haskell $(\lambda x. x)$ apple \Rightarrow apple

Computation by Calculation

Substituting equals by equals

$$\begin{aligned}
 & (2+3) * (5-1) \\
 & \quad \parallel \\
 & = 5 * (5-1) \\
 & \quad \quad \parallel \\
 & = 5 * 4 \\
 & \quad \quad \parallel \\
 & = 20
 \end{aligned}$$

Computation via Substituting Equals by Equals

```
(1 + 3) * (4 + 5)
-- subst 1 + 3 = 4
==> 4 * (4 + 5)
-- subst 4 + 5 = 9
==> 4 * 9
-- subst 4 * 9 = 36
==> 36
```

Computation via Substituting Equals by Equals

Equality-Substitution enables **Abstraction** via **Pattern Recognition**

Abstraction via Pattern Recognition

Repeated Expressions

1. $31 * (42 + 56)$
2. $70 * (12 + 95)$
3. $90 * (68 + 12)$

Recognize Pattern as λ -function

pat = \x y z -> x * (y + z)

Equivalent Haskell Definition

pat x y z = x * (y + z)

Function Call is Pattern Instance

(pat 31 42 56) ==> 31 * (42 + 56) ==> 31 * 98 ==> 3038
 pat 70 12 95 ==> 70 * (12 + 95) ==> 70 * 107 ==> 7490
 pat 90 68 12 ==> 90 * (68 + 12) ==> 90 * 80 ==> 7200

pat x y z = x * (y + z)

pat x y = \z -> x * (y + z)

pat x = \y z -> x * (y + z)

pat = \x y z -> x * (y + z)

Key Idea: Computation is substitute equals by equals.

CSE 30

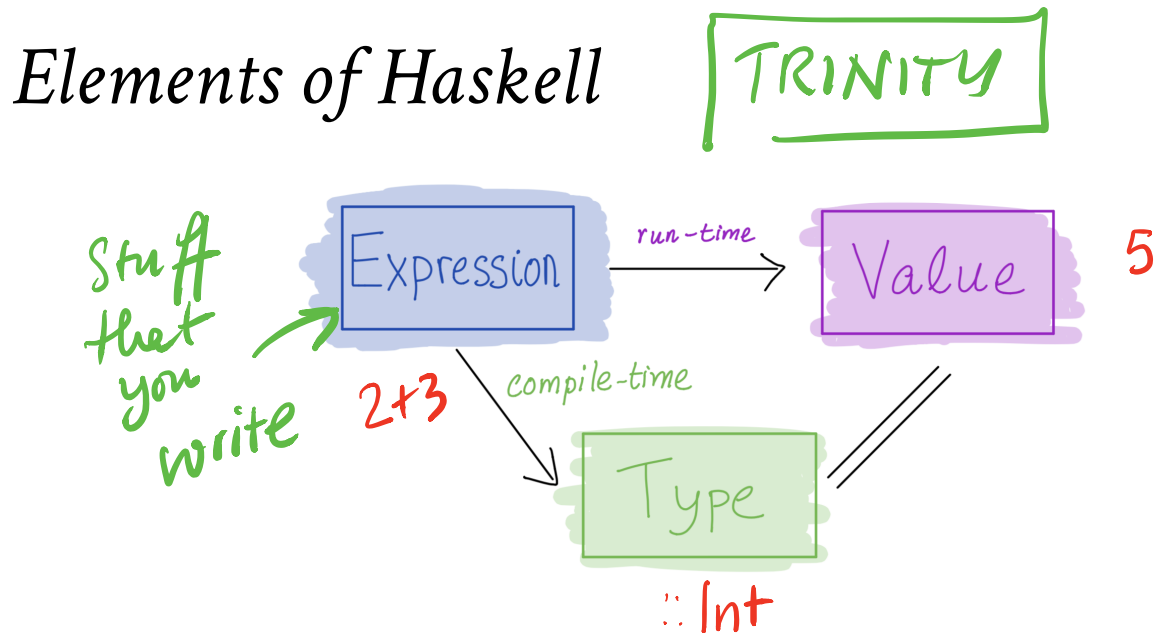
Reg + FRAMES

CSE 131

Programming in Haskell

Substitute Equals by Equals

Thats it! (Do not think of registers, stacks, frames etc.)

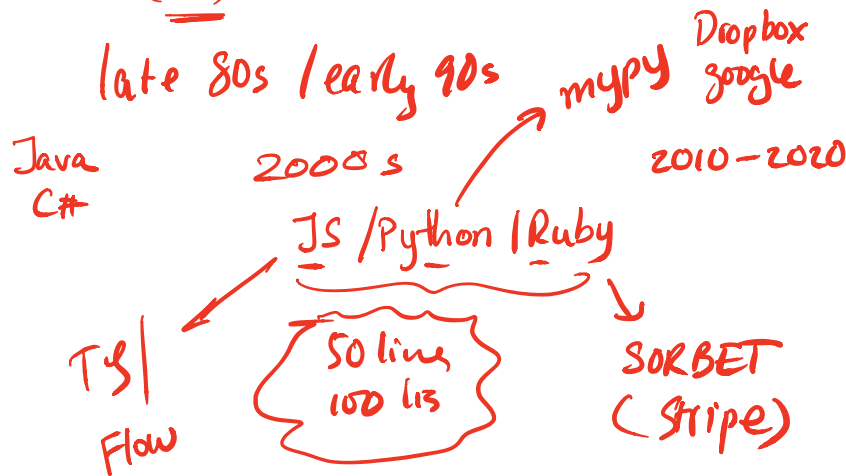


- Core program element is an **expression**
- Every *valid* expression has a **type** (determined at compile-time)
- Every *valid* expression reduces to a *value* (computed at run-time)

Ill-typed* expressions are rejected at *compile-time* before execution

- like in Java (*statically*)
- not like λ -calculus or Python... (*mypy*)

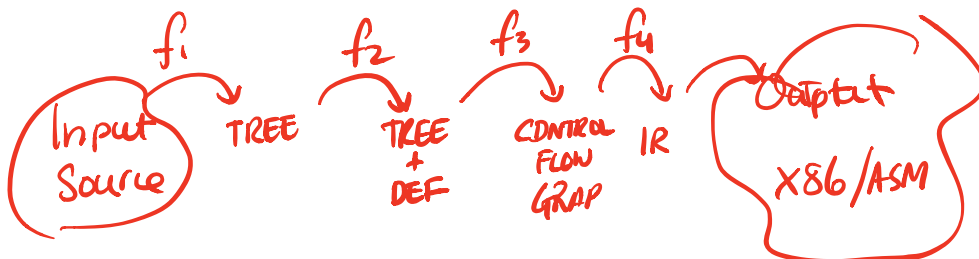
weirdo = (1 0) -- rejected by GHC



$f : \underline{\text{Int}} \rightarrow \underline{\quad}$

Why are types good?

- Helps with program *design*
- Types are contracts (ignore ill-typed inputs!)
- Catches errors *early*
- Allows compiler to *generate code*
- Enables compiler optimizations



01 - haskell

The Haskell Eco-System

- **Batch compiler:** `ghc` Compile and run large programs
- **Interactive Shell** `ghci` Shell to interactively run small programs online (<https://repl.it/languages/haskell>)
- **Build Tool** `stack` Build tool to manage libraries etc.

Interactive Shell: ghci

```
$ stack ghci
```

```
:load file.hs
```

```
:type expression
```

```
:info variable
```

A Haskell Source File

A sequence of **top-level definitions** x_1, x_2, \dots

- Each has *type* $\text{type}_1, \text{type}_2, \dots$
- Each defined by *expression* $\text{expr}_1, \text{expr}_2, \dots$

$x_1 :: \text{type}_1$

$x_1 = \text{expr}_1$

$x_2 :: \text{type}_2$

$x_2 = \text{expr}_2$

•
•
•

 *type anot*

Basic Types

```
ex1 :: Int
```

```
ex1 = 31 * (42 + 56)  -- this is a comment
```

```
ex2 :: Double
```

```
ex2 = 3 * (4.2 + 5.6)  -- arithmetic operators "overloaded"
```

```
ex3 :: Char
```

```
ex3 = 'a'  -- 'a', 'b', 'c', etc. built-in `Char` values
```

```
ex4 :: Bool
```

```
ex4 = True  -- True, False are builtin Bool values
```

```
ex5 :: Bool
```

```
ex5 = False
```

QUIZ: Basic Operations

QUIZ

```
ex6 :: Int
```

```
ex6 = 4 + 5
```

```
ex7 :: Int
```

```
ex7 = 4 * 5
```

```
ex8 :: Bool
```

```
ex8 = 5 > 4
```

```
quiz :: ???
```

```
quiz = if ex8 then ex6 else ex7
```

What is the *type* of quiz?

- A. Int
- B. Bool
- C. Error!

QUIZ: Basic Operations

```
ex6 :: Int
ex6 = 4 + 5
```

```
ex7 :: Int
ex7 = 4 * 5
```

```
ex8 :: Bool
ex8 = 5 > 4
```

```
quiz :: ???
quiz = if ex8 then ex6 else ex7
```

What is the *value* of quiz ?

A. 9

B. 20

C. Other!

$e_1 :: \text{Bool}$ $e_2 :: T$ $e_3 :: T$

if e_1 then e_2 else $e_3 :: T$

ill-typed \equiv has no sensible type

Function Types

In Haskell, a **function** is a **value** that has a type

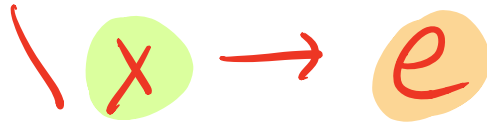
A -> B

A function that

- takes *input* of type A
- returns *output* of type B



For example



```
isPos :: Int -> Bool
```

```
isPos = \n -> (x > 0)
```

Define **function-expressions** using λ like in λ -calculus!

But Haskell also allows us to put the parameter on the *left*

```
isPos :: Int -> Bool
```

```
isPos n = (x > 0)
```

(Meaning is **identical** to above definition with $\lambda n -> \dots$)

Multiple Argument Functions

A function that

- takes three *inputs* A1 , A2 and A3
- returns one *output* B has the type

```
A1 -> A2 -> A3 -> B
```

For example

```
pat :: Int -> Int -> Int -> Int
pat = \x y z -> x * (y + z)
```

which we can write with the params on the *left* as

```
pat :: Int -> Int -> Int -> Int
pat x y z = x * (y + z)
```

QUIZ

What is the type of quiz ?

```
quiz :: ???
quiz x y = (x + y) > 0
```

- A. Int -> Int
- B. Int -> Bool
- C. Int -> Int -> Int
- D. Int -> Int -> Bool
- E. (Int, Int) -> Bool

Function Calls

A function call is *exactly* like in the λ -calculus

func
e1 e2 *arg*

where e1 is a function and e2 is the argument. For example

```
>>> isPos 12
```

```
True
```

```
>>> isPos (0 - 5)
```

```
False
```

Multiple Argument Calls

With multiple arguments, just pass them in one by one, e.g.

```
((e e1) e2) e3)
```

For example

```
>>> pat 31 42 56
3038
```

EXERCISE

Write a function `myMax` that returns the *maximum* of two inputs

```
myMax :: Int -> Int -> Int
myMax = ???
```

When you are done you should see the following behavior:

```
>>> myMax 10 20
20
```

```
>>> myMax 100 5
100
```

EXERCISE

Write a function `sumTo` such that `sumTo n` evaluates to $0 + 1 + 2 + \dots + n$

```
sumTo :: Int -> Int
sumTo n = ???
```

When you are done you should see the following behavior:

```
>>> sumTo 3
6
>>> sumTo 4
10
>>> sumTo 5
15
```