Datatypes and Recursion

$$
\begin{aligned}
& \text { Plan for this week } \\
& \text { int } \\
& \text { boo } \\
& \text { chat, 4-2 } \\
& \text { ( } 2, \text { "cf") } \\
& \text { - writing functions using pattern matching and recursion } \\
& \text { This week: } \\
& {[2,3,4]}
\end{aligned}
$$



- user-defined data types
- and how to manipulate them using pattern matching and recursion
- more details about recursion


## Representing complex data

We've seen:

- base types: Sol, Int, Integer, Float
- some ways to build up types: given types T1, T2
- functions: T1 -> T2
- tuples: (T1, T2)
- lists: [T1]

$$
\begin{aligned}
& T=\ln t \mid \text { Bool|Char } \\
& 1(T, T) \\
& 1\left(T_{1}, \ldots, T_{k}\right) \\
& |[T]| T \rightarrow T
\end{aligned}
$$

Algebraic Data Types: a single, powerful technique for building up types to represent complex data

- Lets you define your own data types
- Tuples and lists are special cases


## Building data types

(Int)
string)
Three key ways to build complex types/values:

1. Product types (each-of): a value of $T$ contains a value of $T$ and value of $T 2$
2. Sum types (one-of): a value of T contains a value of T or a value of T 2 "Union"
3. Recursive types: a value of $T$ contains a sub-value of the same type $T$

```
deadlineDate :: (Int, Int, Int)
deadlineDate = (2, 4, 2019)
deadlineTime :: (Int, Int, Int)
deadlineTime = (11, 59, 59)
-- / Deadline date extended by one day
extension :: (Int, Int, Int) -> (Int, Int, Int)
extension = ...
```

Can you spot them?

## 1. Verbose and unreadable

A type synonym for $T$ : a name that can be used interchangeably with $T$

```
type Date = (Int, Int, Int)
type Time = (Int, Int, Int)
deadlineDate :: Date
deadlineDate = (2, 4, 2019)
deadlineTime :: Time
deadlineTime = (11, 59, 59)
-- / Deadline date extended by one day
extension :: Date -> Date
extension = ...
```


## 2. Unsafe

We want this to fail at compile time!!!
extension deadlineTime

Solution: construct two different datatypes

```
data Date = Date Int Int Int
data Time = Time Int Int Int
-- constructor^ ^parameter types
deadlineDate :: Date
deadlineDate = Date 2 4 2019
deadlineTime :: Time
deadlineTime = Time 11 59 59
```


## Record syntax

Haskell's record syntax allows you to name the constructor parameters:

- Instead of
data Date = Date Int Int Int
- you can write:


$$
\begin{aligned}
& \text { MkDate:: |ut } \rightarrow \text { Mut }\left.\rightarrow\right|_{u} \rightarrow \text { Dat } \\
& \text { Destr. } \\
& \text { month: Dat } \rightarrow \text { Mut }
\end{aligned}
$$

- then you can do:
deadlineDate = Date 242019
dealineMonth $=$ month deadlineDate - - yikes, use field name as a function


## Building data types



Three key ways to build complex types/values:

1. Product types (each-of): a value of T contains a value of T1 and a value of T2 [done]
2. Sum types (one-of): a value of T contains a value of T1 or a value of T2
3. Recursive types: a value of $T$ contains a sub-value of the same type $T$

## Example: NanoMarkdown

Suppose I want to represent a text document with simple markup
Each paragraph is either:

- plain text (String)
$\rightarrow$ - heading: level and text (Int and String)
$\rightarrow$ list: ordered? and items (Bool and [String])
$\langle 0 l\rangle-. .-\langle 101\rangle$
I want to store all paragraphs in a list
$\langle u l\rangle \ldots\langle i u l\rangle$
doc = [ (1, "Notes from 130")
-- Lvl 1 heading
, "There are two types of languages:"
-- Plain text
, (True, ["those people complain about", "those no one uses"]) -- Ordered list ]

But this does not type check!!!

## Sum Types

Solution: construct a new type for paragraphs that is a sum (one-of) the three options!
Each paragraph is either:

- plain text (String)
- heading: level and text ( Int and String)
- list: ordered? and items (Bool and [String])

```
data Paragraph -- ^ 3 constructors, w/ different parameters
    = PText String -- ^ text : plain string
    | PHeading Int String -- ^ heading: level and text (`Int` and `String`)
    | PList Bool [String] -- ^ list : ordered? and items (`Bool` and `[String]`)
```


## QUIZ

data Paragraph
= PText String
| PLeading Int String
| PList Boot [String]
(P Text "hey")
What is the type of(PText "Hey there! ") ? ie. How would GHCi reply to:
>:t (PText "Hey there!")
A. Syntax error
B. Type error
C. PText
D. String
E. Paragraph


## Constructing datatypes



- C1 .. Cn are the constructors of $T$

A value of type $T$ is

- either C1 v1 .. vk with vi :: T1i
- or C2 v1 .. vl with vi :: T2i
- or ...
- or Cn v1 .. vm with vi : : Tni

You can think of a $T$ value as a box:

- either a box labeled C1 with values of types T11 .. T1k inside
- or a box labeled C2 with values of types T21 .. T2l inside
- or ...
- or a box labeled Cn with values of types Tn 1 .. Tnm inside


One-of Types

Apply a constructor = pack some values into a box (and label it)

- PText "Hey there!"
o put "Hey there!" in a box labeled PText
- PHeading 1 "Introduction"
- put 1 and "Introduction" in a box labeled PHeading
- Boxes have different labels but same type ( Рaragraph )


The Paragraph Type with example values:


The Paragraph Type

## QUIZ

data Paragraph
= PText String
| PHeading Int String
| PList Bool [String]

What would GHCi say to
>:t [PHeading 1 "Introduction", Fext "Hey there!"]
A. Syntax error
B. Type error
C. Paragraph
D. [Paragraph]
E. [String]

## Example: NanoMD

## data Рагаgгарh

= PText String
| PHeading Int String
| PList Bool [String]
Now I can create a document like so:

```
doc :: [Рагаgraph]
doc = [ PHeading 1 "Notes from 130"
    , PText "There are two types of languages:"
    , PList True ["those people complain about", "those no one uses"])
    ]
```

Now I want convert documents in to HTML.
I need to write a function:
html :: Рaragraph -> String
html p = ??? -- depends on the kind of paragraph!

How to tell what's in the box?

- Look at the label!


## Pattern matching

Pattern matching = looking at the label and extracting values from the box

- we've seen it before
- but now for arbitrary datatypes
html :: Paragraph -> String
html (PText str) int= s.fr -- It's a plain text! Get string
html (PLeading lvi str) $\xlongequal[=]{\boldsymbol{Z}} . .$. .- It's a heading! Get level and string
html (PList ord $\overline{i t e m s})=\ldots$ - It's a list! Get ordered and items


```
html :: Paragraph -> String
html (PText str) -- It's a plain text! Get string
    = unlines [open "p", str, close "p"]
html (PHeading lvl str) -- It's a heading! Get level and string
    = let htag = "h" ++ show lvl
    in unwords [open htag, str, close htag]
html (PList ord items) -- It's a list! Get ordered and items
    = let ltag = if ord then "ol" else "ul"
    litems = [unwords [open "li", i, close "li"] | i <- items]
    in unlines ([open ltag] ++ litems ++ [close ltag])
```


## Dangers of pattern matching (1)

```
html :: Paragraph -> String
html (PText str) = ...
html (PList ord items) = ...
```

What would GHCi say to:

```
html (PHeading 1 "Introduction")
```


## Dangers of pattern matching (2)

```
html :: Paragraph -> String
html (PText str) = unlines [open "p", str, close "p"]
html (PHeading lvl str) = ...
html (PHeading 0 str) = html (PHeading 1 str)
html (PList ord items) = ...
```

What would GHCi say to:
html (PHeading 0 "Introduction")

## Dangers of pattern matching

Beware of missing and overlapped patterns

- GHC warns you about overlapped patterns
- GHC warns you about missing patterns when called with -W (yse : set -W in GHCi)


## Pattern-Match Expression

Everything is an expression?


We've seen: pattern matching in equations
Actually, pattern-match is also an expression


The code we saw earlier was syntactic sugar
html $($ C1 $\times 1 \ldots)=e 1$
html (C2 x2 ...) $=$ e2
html (C3 x3 ...) $=$ e3
is just for humans, internally represented as a case-of expression
html $p=$ case $p$ of
(C1 x1 ...) -> e1
(C2 x2 ...) -> e2
(C3 x3 ...) -> e3

## QUIZ

What is the type of
case Text "Hey" of
Text str $\rightarrow$ str
PHead lvi $\rightarrow$ elul
Plist ord $-\rightarrow$ ord
B. Type error
$\tau$
C String_
D. Paragraph
E. Paragraph -> String

## Pattern matching expression: typing

The case expression

```
case e of
    pattern1 -> e1
    pattern2 -> e2
    patternN -> eN
```

has type T if

- each e1 ... eN has type T
- e has some type D
- each pattern1 ... patternN is a valid pattern for D
- i.e. a variable or a constructor of $D$ applied to other patterns

The expression e is called the match scrutinee

## QUIZ

What is the type of
let $p=$ Text "Hey there!"
in case $p$ of
PText _ -> 1
PHeading _ _ -> 2
PList _ _ -> 3
A. Syntax error
B. Type error
C. Paragraph
D. Int
E. Paragraph -> Int

## Building data types



