

# CSE 130 Midterm Solution, Spring 2018

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## Part I. Lambda Calculus [60 pts]

### Q1: Reductions [20 pts]

#### 1.1 [5 pts]

$(\lambda x \rightarrow x (\lambda x \rightarrow x)) (f x)$

- (A)  $=b> f x (\lambda x \rightarrow x)$  valid
- (B)  $=b> f x (\lambda x \rightarrow f x)$
- (C)  $=b> f x (\lambda f x \rightarrow f x)$
- (D)  $=a> (\lambda y \rightarrow y (\lambda x \rightarrow x)) (f y)$
- (E)  $=a> (\lambda x \rightarrow x (\lambda y \rightarrow y)) (f x)$  valid

#### 1.2 [5 pts]

$\lambda x \rightarrow (\lambda y z \rightarrow x y) (\lambda x \rightarrow x)$

- (A)  $=a> \lambda x \rightarrow (\lambda y x \rightarrow x x) (\lambda x \rightarrow x)$
- (B)  $=a> \lambda a \rightarrow (\lambda y z \rightarrow a y) (\lambda x \rightarrow a)$
- (C)  $=*> \lambda a \rightarrow (\lambda y z \rightarrow a y) (\lambda a \rightarrow a)$  valid
- (D)  $=b> \lambda x z \rightarrow x (\lambda x \rightarrow x)$  valid

(E)  $=b> \lambda y z \rightarrow (\lambda x \rightarrow x) y$

### 1.3 [5 pts]

$(\lambda f g x \rightarrow f (g x)) (\lambda x \rightarrow g x) (\lambda z \rightarrow z)$

(A)  $=b> (\lambda f g x \rightarrow f (g x)) (g (\lambda z \rightarrow z))$

(B)  $=b> (\lambda g x \rightarrow (\lambda x \rightarrow g x) (g x)) (\lambda z \rightarrow z)$

(C)  $=*> \lambda x \rightarrow g x$  valid

(D)  $=*> \lambda y \rightarrow g y$  valid

(E)  $=*> \lambda x \rightarrow f x$

### 1.4 [5 pts]

$(\lambda x y \rightarrow \lambda b u v \rightarrow b v u) (\lambda x y \rightarrow y) (\lambda x y \rightarrow y) (\lambda x y \rightarrow x)$

(A)  $=b> (\lambda y \rightarrow \lambda b u v \rightarrow b v u) (\lambda x y \rightarrow y) (\lambda x y \rightarrow x)$  valid

(B)  $=b> (\lambda y \rightarrow \lambda b u v \rightarrow b v u) (\lambda x y \rightarrow y) (\lambda x y \rightarrow y)$

(C)  $=*> (\lambda b u v \rightarrow b v u) (\lambda x y \rightarrow x)$  valid

(D)  $=*> \lambda x y \rightarrow y$  valid

(E)  $=b> \lambda y \rightarrow (\lambda x y \rightarrow y) (\lambda x y \rightarrow y) (\lambda x y \rightarrow x)$

## Q2: Lists [40 pts]

### 2.1 Repeat [10 pts]

```
let REPEAT = \n x -> n (PAIR x) FALSE
```

### 2.2 Empty\* [20 pts]

```
let EMPTY = \xs -> xs (\x y z -> FALSE) TRUE
```

Alternatively:

```
let EMPTY = \xs -> xs (\x y -> NOT) TRUE
```

### 2.3 Length [10 pts]

```
let LEN = FIX (\rec l -> ITE (EMPTY l) ZERO (INC (rec (SND l))))
```

## Part II. Datatypes and Recursion [50 pts]

### Q3: Binary Search Trees [50 pts]

#### 3.1 Size [5 pts]

```
size :: Tree -> Int
size Empty = 0
size (Node _ l r) = 1 + size l + size r
```

#### 3.2 Insert [10 pts]

```
insert :: Int -> Tree -> Tree
insert x Empty = Node x Empty Empty
insert x (Node y l r)
| x == y    = Node y l r
| x < y     = Node y (insert x l) r
| otherwise  = Node y l (insert x r)
```

#### 3.3 Sort [15 pts]

```
sort :: [Int] -> [Int]
sort xs = toList (fromList xs)
where
  fromList :: [Int] -> Tree
  fromList [] = Empty
```

```

fromList (x:xs) = insert x (fromList xs)

toList :: Tree -> [Int]
toList Empty = []
toList (Node x l r) = toList l ++ [x] ++ toList r

```

### 3.4 Tail-recursive size\* [20 pts]

```

size :: Tree -> Int
size t = loop 0 [] t
  where
    loop :: Int -> [Tree] -> Tree -> Int
    loop acc []     Empty      = acc
    loop acc (t:ts) Empty      = loop acc      ts      t
    loop acc ts     (Node _ l r) = loop (acc + 1) (r:ts) l

```

Alternatively:

```

size :: Tree -> Int
size t = loop 0 Empty t
  where
    loop :: Int -> Tree -> Tree -> Int
    loop acc Empty      Empty      = acc
    loop acc (Node _ l r) Empty      = loop (acc + 1) l      r
    loop acc t           (Node x l r) = loop acc      (Node x t r) l

```