Higher-Order Functions

Plan for this week

Last week:

- user-defined data types
- manipulating data-types with pattern matching and recursion

Calendar

Contact

Grades

Lectures

Assignments

Links

Piazza

Canvas

how to make recursive functions more efficient with *tail recursion*

The long arc of history

Pattern matching is a very old PL idea ...

- Variants of LISP from 1970 by Fred McBride
- ... but will finally be added to Python 3.10 (2020)
- https://www.python.org/dev/peps/pep-0622/
- def make_point_3d(pt):
 match pt:
 case (x, y):
 return Point3d(x, y, 0)
 case (x, y, z):
 return Point3d(x, y, z)
 case Point2d(x, y):
 return Point3d(x, y, 0)
 case Point3d(_, _, _):
 return pt
 case _:
 raise TypeError("not a point we support")

Plan for this week

Last week:

- user-defined data types
- manipulating data-types with pattern matching and recursion
- how to make recursive functions more efficient with tail recursion
- This week:
 - code reuse with higher-order functions (HOFs)
 - \bullet some useful HOFs: <code>map</code> , <code>filter</code> , and <code>fold</code>

Recursion is good...

- Recursive code mirrors recursive data
 - Base constructor -> Base case
 - $\circ~$ Inductive constructor -> Inductive case (with recursive call)
- But it can get kinda repetitive!



Example: four-letter words

Let's write a function fourChars :

-- fourChars [] ==> [] -- fourChars ["i", "must", "do", "work"] ==> ["must", "work"]

fourChars :: [String] -> [String]
fourChars [] = ...
fourChars (x:xs) = ...

Yikes! Most Code is the Same!

Lets rename the functions to foo:

foo [] = []
foo (x:xs)
 | x mod 2 == 0 = x : foo xs
 | otherwise = foo xs
foo [] = []
foo (x:xs)
 | length x == 4 = x : foo xs
 | otherwise = foo xs

Only difference is **condition**

• x mod 2 == 0 vs length x == 4

Moral of the day

D.R.Y. Don't Repeat Yourself!

Can we

- reuse the general pattern and
- *plug-in* the custom condition?

Higher-Order Functions

General **Pattern**

- expressed as a higher-order function
- takes plugin operations as arguments

Specific Operation

• passed in as an argument to the HOF

The "filter" pattern



The filter Pattern

General Pattern

• HOF filter

Recursively traverse list and pick out elements that satisfy a predicate

Specific Operations

• Predicates isEven and isFour

| <pre>filter f [] filter f (x:xs)</pre> | = [] |
|--|---|
| <mark>f x</mark> otherwise | = x : filter f xs = filter f xs |
| <pre>evens = filter isEven where isEven x = x `mod` 2 == 0</pre> | <pre>fourChars = filter isFour where isFour x = length x == 4</pre> |
| filter instances | |

Avoid duplicating code!

QUIZ: What is the type of *filter*?

evens :: [Int] -> [Int] evens xs filter isEven xs where isEven :: Int -> Bool

| is Even $x = x \mod 2 = 0$ |
|--|
| <pre> fourChars ["i", "must", "do", "work"] ==> ["must", "work"] fourChars :: [String] -> [String] fourChars xs = filter isFour xs where isFour :: String -> Bool isFour = leasth w are 4</pre> |
| So what's the type of filter? |
| <pre>{- A -} filter :: (Int -> Bool) -> [Int] -> [Int] {- B -} filter :: (String -> Bool) -> [String] -> [String]</pre> |
| {- C -} filter :: (a -> Bool) -> [a] -> [a] |
| <i>{- D -}</i> filter :: (a -> Bool) -> [a] -> [Bool] <i>{- E -}</i> filter :: (a -> b) -> [a] -> [b] |

Type of *filter*

-- evens [1,2,3,4] ==> [2,4] evens :: [Int] -> [Int] evens xs = filter isEven xs where isEven :: Int -> Bool isEven x = x `mod` 2 == 0 -- fourChars ["i", "must", "do", "work"] ==> ["must", "work"] fourChars :: [String] -> [String] fourChars xs = filter isFour xs where isFour :: String -> Bool isFour x = length x == 4For any type a • Input a predicate a -> Bool and collection [a] • Output a (smaller) collection [a] filter :: (a -> Bool) -> [a] -> [a] filter does not care what the list elements are • as long as the predicate can handle them

filter is **polymorphic** (generic) in the type of list elements

Example: ALL CAPS!

Lets write a function shout :

-- shout [] ==> [] -- shout ['h','e','l','l','o'] ==> ['H','E','L','L','0']

shout :: [Char] -> [Char]
shout [] = ...
shout (x:xs) = ...

Example: squares

Lets write a function squares :

-- squares [] ==> [] -- squares [1,2,3,4] ==> [1,4,9,16]

squares :: [Int] -> [Int]
squares [] = ...
squares (x:xs) = ...

Yikes, Most Code is the Same

```
-- shout
foo [] = []
foo (x:xs) = toUpper x : foo xs
-- squares
foo [] = []
foo (x:xs) = (x * x) : foo xs
```

Lets **refactor** into the **common pattern**

pattern = ...

The "map" pattern

shout [] = [] squares [] = [] shout (x:xs) = toUpper x : shout xs squares (x:xs) = (x*x) : squares xs map <mark>f</mark> [] = [] map f (x:xs) = f x : map f xs The map Pattern **General** Pattern • HOF map • Apply a transformation f to each element of a list Specific Operations • Transformations to Upper and $x \rightarrow x * x$ map f [] = [] map f (x:xs) = f x : map f xs

map f [] = [] map f (x:xs) = f x : map f xs $shout = map (\x -> toUpper x)$ $squares = map (\x -> x*x)$ map instances

QUIZ ords :: [Char] -> [hut] ords K = map ord KS What is the type of map? map f [] = [] map f (x:xs) = f x : map f xS (A) (Char -> Char) -> [Char] -> [Char] X (B) (Int -> Int) -> [Int] -> [Int] X (C) (a -> a) -> [a] -> [a] Z (D) (a -> b) -> [a] -> [b] Z (E) (a -> b) -> [c] -> [d]

-- For any types `a` and `b`
-- if you give me a transformation from `a` to `b`
-- and a list of `a`s,
-- I'll give you back a list of `b`s
map :: (a -> b) -> [a] -> [b]

Type says it all!

- The only meaningful thing a function of this type can do is apply its first argument to elements of the list
- Hoogle it!

Things to try at home:

- can you write a function map' :: (a -> b) -> [a] -> [b] whose behavior is different from map?
- can you write a function map' :: (a -> b) -> [a] -> [b] such that map' f xs returns a list whose elements are not in map f xs?



Don't Repeat Yourself

- Benefits of factoring code with HOFs:Reuse iteration pattern
 - think in terms of standard patterns
 - ∘ less to write
 - easier to communicate
 - Avoid bugs due to repetition

Recall: length of a list

-- len [] ==> 0 -- len ["carne", "asada"] ==> 2 len :: [a] -> Int len [] = 0 len (x:xs) = 1 + len xs

Recall: summing a list

-- sum [] ==> 0 -- sum [1,2,3] ==> 6 sum :: [Int] -> Int sum [] = 0 sum (x:xs) = x + sum xs

Example: string concatenation

Let's write a function cat: -- cat [] ==> "" -- cat ["carne", "asada", "torta"] ==> "carneasadatorta" cat :: [String] -> String cat [] = ... cat (x:xs) = ...

Can you spot the pattern?

-- len foo [] = 0 foo (x:xs) = 1 + foo xs

-- sum foo [] = 0 foo (x:xs) = x + foo xs

-- cat foo [] = "" foo (x:xs) = x ++ foo xs

pattern = ...

The "fold-right" pattern $\begin{bmatrix} len [] &= 0 \\ len (x:xs) &= 1 + len xs \end{bmatrix} \begin{bmatrix} sum [] &= 0 \\ sum (x:xs) &= x + sum xs \end{bmatrix} \begin{bmatrix} cat [] &= "" \\ cat (x:xs) &= x + t + sum xs \end{bmatrix}$ $\begin{bmatrix} foldr f b [] &= b \\ foldr f b (x:xs) &= f x (foldr f b xs) \end{bmatrix}$

The foldr Pattern

General Pattern

Recurse on tailCombine result with the head using some binary operation

foldr f b [] = b foldr f b (x:xs) = f x (foldr f b xs)

Let's refactor sum, len and cat:

sum = foldr

cat = foldr

len = foldr

Factor the recursion out! $\begin{array}{c} (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b \\ (T_x \rightarrow T_b \rightarrow T_b) \rightarrow T_b \rightarrow [T_x] \rightarrow T_b \end{array}$

 $\begin{array}{c} \chi :: T_{\chi} \\ \chi :: [T_{\chi}] \\ f :: T_{\chi} \to T_{b} \longrightarrow T_{b} \\ \hline T_{f} \end{array}$

foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs) $T_{r} T_{b}$ len = foldr (\x n -> 1 + n) 0
sum = foldr (\x n -> x + n) 0
cat = foldr (\x s -> x ++ n) ""

foldr instances

You can write it more clearly as

sum = foldr (+) 0

cat = foldr (++) ""

The "fold-right" pattern (a3 6))) a, 'o' (a2 `o' ۰0 (au 6 foldr f b [a1, a2, a3, a4] ==> f a1 (foldr f b [a2, a3, a4]) ==> f a1 (f a2 (foldr f b [a3, a4])) ==> f a1 (f a2 (f a3 (foldr f b [a4]))) ==> f a1 (f a2 (f a3 (f a4 (foldr f b [])))) ==> f a1 (f a2 (f a3 (f a4 b))) Accumulate the values from the right 11 For example: D foldr (+) 0 [1, 2, 3, 4] ./// ==> 1 + (foldr (+) 0 [2, 3, 4]) ==> 1 + (2 + (foldr (+) 0 [3, 4])) ==> 1 + (2 + (3 + (foldr (+) 0 [4]))) ==> 1 + (2 + (3 + (4 + (foldr (+) 0 [])))) ==> 1 + (2 + (3 + (4 + 0)))

QUIZ What does this evaluate to? foldr f b [] = b foldr f b (x:xs) = f x (foldr f b xs) quiz = foldr (\x v -> x : v) [] [1,2,3]

11

b op (A) Type error 1 'op' (foldr o [] (2:3:[])) **(B)** [1,2,3] 'op' (2 'o' (Gibr o E3 (3:E3))) (C) [3,2,1] (z'op' (3'op' (E])) -> 1 °op` (D) [[3],[2],[1]] (E) [[1],[2],[3]]

foldr (:) [] [1,2,3] ==> (:) 1 (foldr (:) [] [2, 3]) ==> (:) 1 ((:) 2 (foldr (:) [] [3])) ==> (:) 1 ((:) 2 ((:) 3 (foldr (:) [] []))) ==> (:) 1 ((:) 2 ((:) 3 [])) == 1 : (2 : (3 : [])) == [1,2,3] len E] = 0 leu (x:xS) = 1 + lu xS lew = foldr $\underbrace{3}_{(l+)}_{(1x \to l+x)}$

QUIZ

What is the most general type of foldr?

foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)
(A) (a -> a -> a) -> a -> [a] -> a
(B) (a -> a -> b) -> a -> [a] -> b
(C) (a -> b -> a) -> b -> [a] -> b
(D) (a -> b -> b) -> b -> [a] -> b
(E) (b -> a -> b) -> b -> [a] -> b

Tail Recursive Fold

foldr f b [] = b
foldr f b (x:xs) = f x (foldr f b xs)
Is foldr tail recursive?

What about tail-recursive versions?

Let's write tail-recursive sum ! sumTR :: [Int] -> Int sumTR = ...

Lets run ${\tt sumTR}$ to see how it works

| sumTR | [1,2,3] | | | | | | |
|-------|----------|---------|-------|---|---|-----|---|
| ==> | helper 0 | [1,2,3] | | | | | |
| ==> | helper 1 | [2,3] | 0 | + | 1 | ==> | 1 |
| ==> | helper 3 | [3] | 1 | + | 2 | ==> | 3 |
| ==> | helper 6 | [] | 3 | + | 3 | ==> | 6 |
| ==> | 6 | | | | | | |

Note: helper directly returns the result of recursive call!

Let's write tail-recursive cat!
catTR :: [String] -> String
catTR = ...

Lets run catTR to see how it works

| catTR | | ["carne", | "asada", | "torta"] |
|-------------|------------|-----------|----------|----------|
| ==> helper | | ["carne", | "asada", | "torta"] |
| ==> helper | "carne" | [| "asada", | "torta"] |
| ==> helper | "carneasa | ada" | I | "torta"] |
| ==> helper | "carneasa | adatorta" | | [] |
| ==> "carnea | asadatorta | а" | | |

Note: helper directly returns the result of recursive call!

Can you spot the pattern?

| sumTR | | | | | |
|--------|-----|--------|---|--------|---------------|
| foo xs | | | = | helper | 0 xs |
| where | | | | | |
| helper | acc | [] | = | acc | |
| helper | acc | (x:xs) | = | helper | (acc + x) xs |
| | | | | | |
| | | | | | |
| catTR | | | | | |
| foo xs | | | = | helper | "" xs |
| where | | | | | |
| helper | acc | [] | = | асс | |
| helper | acc | (x:xs) | = | helper | (acc ++ x) xs |
| | | | | | |
| | | | | | |

pattern = ...

The "fold-left" pattern



• Use a helper function with an extra accumulator argument

• To compute new accumulator, combine current accumulator with the head using some binary operation

```
foldl f b xs = helper b xs
where
helper acc [] = acc
helper acc (x:xs) = helper (f acc x) xs
```

Let's refactor sumTR and catTR:

sumTR = foldl
catTR = foldl
Factor the tail-recursion out!

QUIZ

What does this evaluate to?
foldl f b xs = helper b xs
where
helper acc [] = acc
helper acc (x:xs) = helper (f acc x) xs
quiz = foldl (\xs x -> x : xs) [] [1,2,3]

(A) Type error
(B) [1,2,3]
(C) [3,2,1]
(D) [[3],[2],[1]]

(E) [[1],[2],[3]]

```
foldl f b (x1: x2: x3 : [])
 ==> helper b (x1: x2: x3 : [])
 ==> helper (f x1 b) (x2: x3 : [])
 ==> helper (f x2 (f x1 b)) (x3 : [])
 ==> helper (f x3 (f x2 (f x1 b))) []
 ==> ( x3 : (x2 : (x1 : [])))
```

The "fold-left" pattern

| foldl | fb | [x1, x2, x3, x4] | | |
|-------|-------------------------|------------------|--|--|
| ==> | helper b | [x1, x2, x3, x4] | | |
| ==> | helper (f b x1) | [x2, x3, x4] | | |
| ==> | helper (f (f b x1) x2) | [x3, x4] | | |
| ==> | helper (f (f (f b x1) x | 2) x3) [x4] | | |
| ==> | helper (f (f (f b x1 |) x2) x3) x4) [] | | |
| ==> | (f (f (f (f b x1) x2) x | 3) x4) | | |
| | | | | |

Accumulate the values from the left

For example:

| foldl | (+) 0 | | [1, 2, 3, 4] |
|-------|---------|--|--------------|
| ==> | helper | 0 | [1, 2, 3, 4] |
| ==> | helper | (0 + 1) | [2, 3, 4] |
| ==> | helper | ((0 + 1) + 2) | [3, 4] |
| ==> | helper | (((0 + 1) + 2)) | + 3) [4] |
| ==> | helper | ((((0 + 1) + 2)) | + 3) + 4) [] |
| ==> | ((((0 + | + 1) + 2) + 3) + | 4) |
| | 5 | and division of the state of th |) |

Left vs. Right

foldl f b [x1, x2, x3] ==> f (f (f b x1) x2) x3 -- Left foldr f b [x1, x2, x3] ==> f x1 (f x2 (f x3 b)) -- *Right* For example: foldl (+) 0 [1, 2, 3] ==> ((0 + 1) + 2) + 3 -- Left foldr (+) 0 [1, 2, 3] ==> 1 + (2 + (3 + 0)) -- Right Different types! foldl :: (b -> a -> b) -> b -> [a] -> b -- Left foldr :: (a -> b -> b) -> b -> [a] -> b -- Right

Higher Order Functions (HOF)

Iteration patterns over collections:

1 OP'

- a->bool • Filter values in a collection given a predicate
- Map (iterate) a given transformation over a collection $a \rightarrow b$
- Fold (reduce) a collection into a value, given a binary operation to combine results (a→b→b)

HOFs can be put into libraries to enable modularity

- Data structure library implements map, filter, fold for its collections
 - generic efficient implementation
 - o generic optimizations: map f (map g xs) --> map (f.g) xs
- Data structure clients use HOFs with specific operations
 - no need to know the implementation of the collection

Crucial foundation of

Same

- "big data" revolution e.g. MapReduce, Spark, TensorFlow
- "web programming" revolution e.g. Jquery, Angular, React

Generated by Hakyll, template by Armin Ronacher, suggest improvements here.